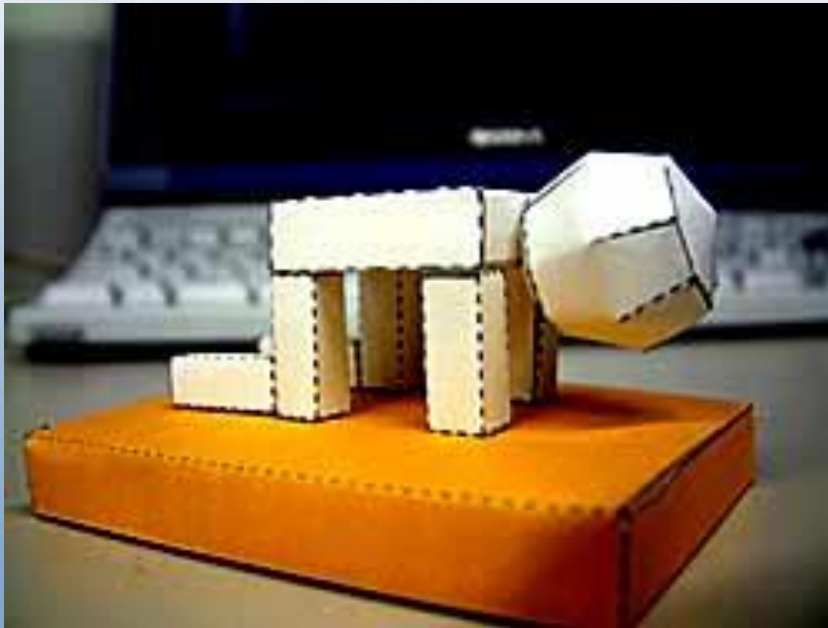
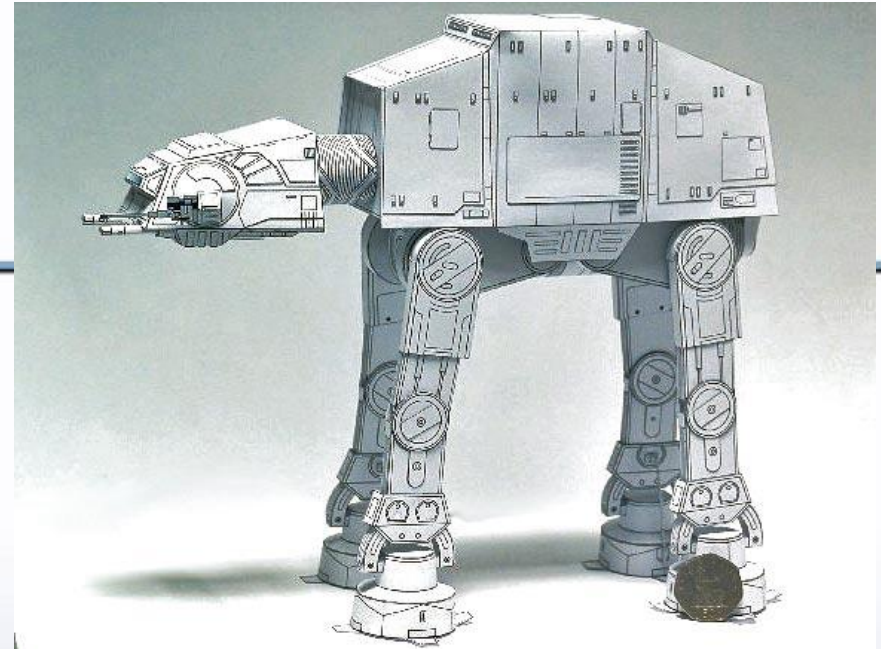




# *BRep and CSG*





# Contents

---

- Boundary Representation BRep
- Euler's formula
- Winged-edge structure
- Euler operators
- CSG tree
- Point-membership classification algorithm
- References



# BRep notions

---

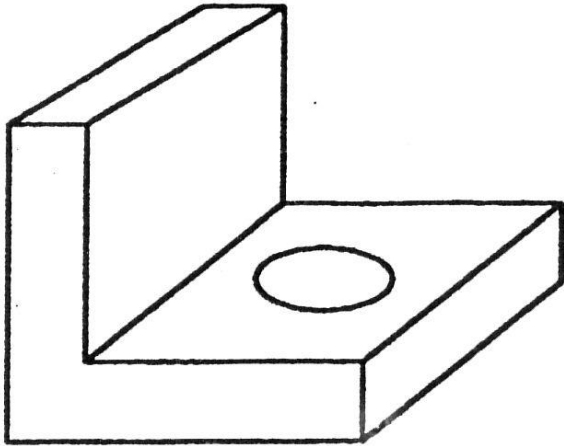
In the **Boundary Representation BRep**, a solid is represented by segmenting its boundary into a finite number of bounded subsets usually called “*faces*” or “*patches*”, and representing each face by its bounding *edges* and *vertices*.



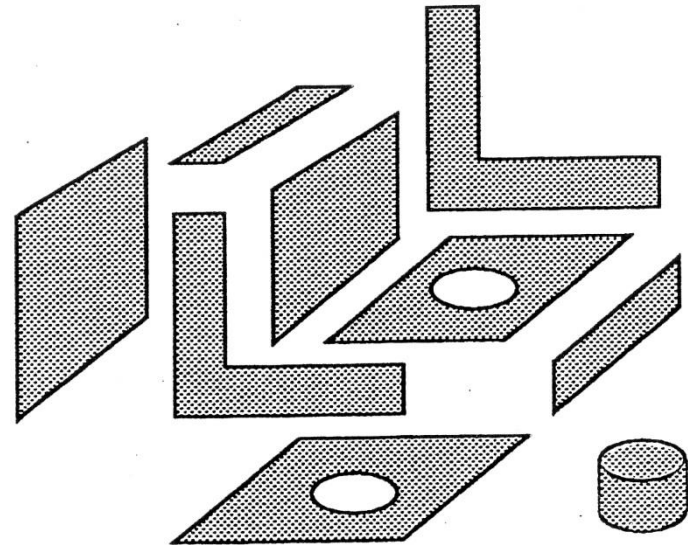
- This description has two parts, a **topological** description of the connectivity and orientation of vertices, edges, and faces in the form of a **graph**, and a **geometric** description for embedding these surface elements in space.
- The topological description specifies vertices, edges, and faces abstractly, and indicates their incidences and adjacencies.
- The geometric description specifies, for example, the coordinates of vertices or the equations of the surfaces containing the faces.



# BRep example



Solid

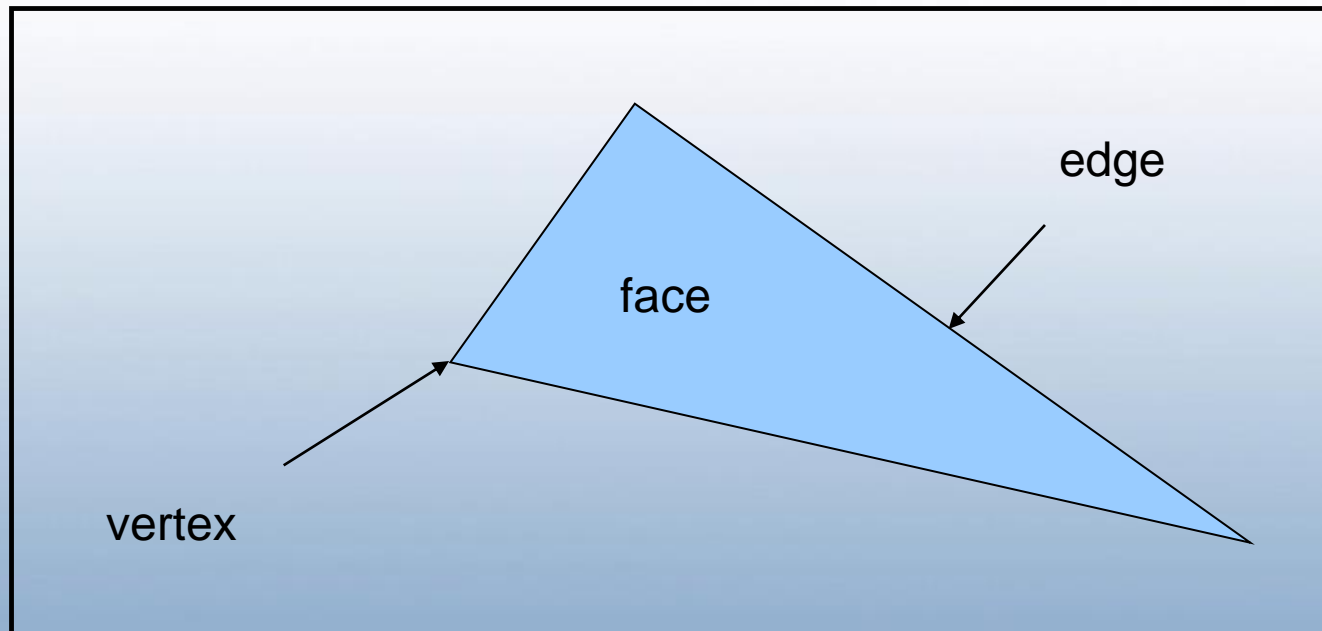


BRep



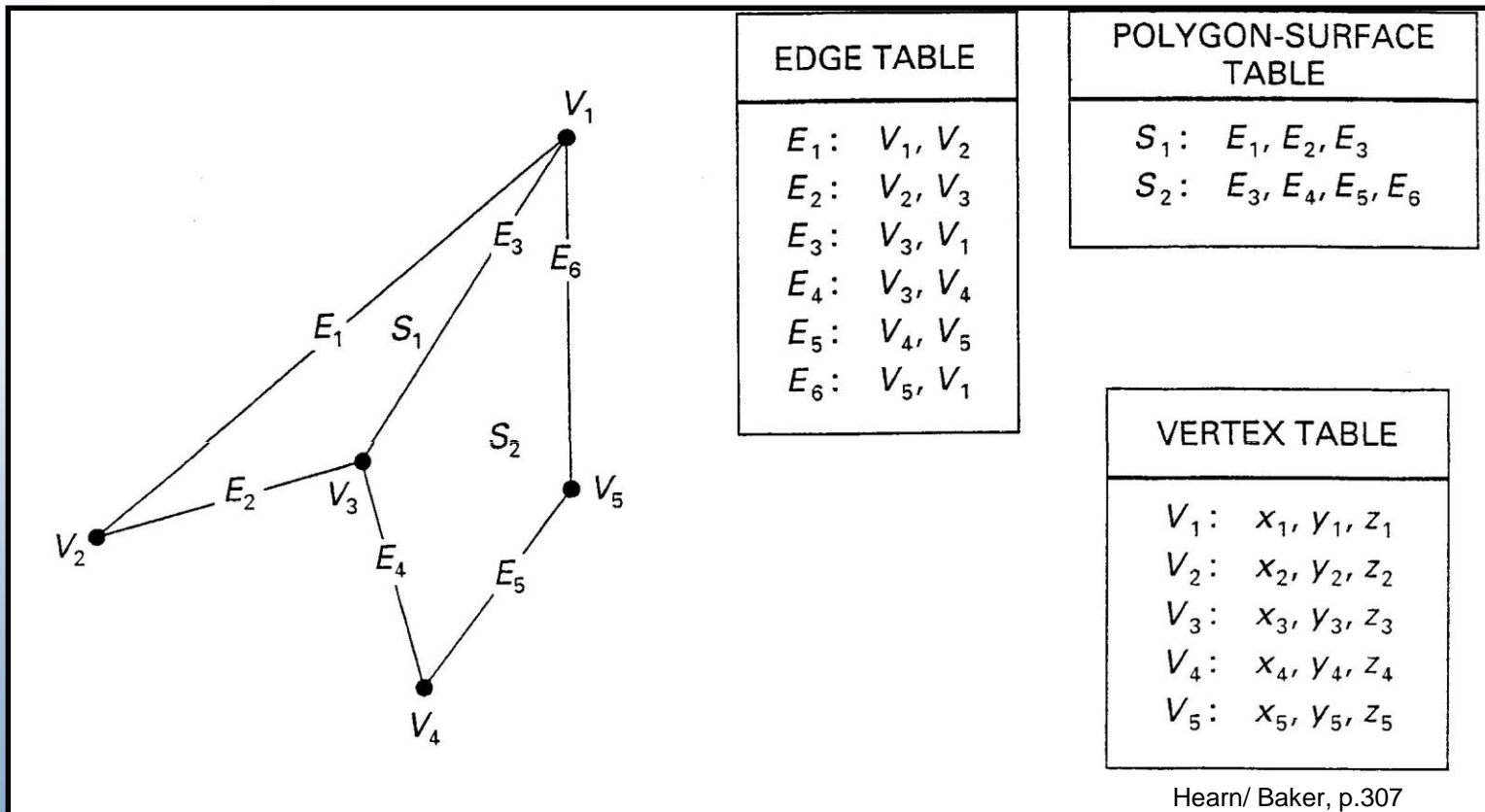
# BRep structure

Historically, BRep evolved from a description of polyhedra in computer graphics:





# BRep structure





# BRep and Graphs

---

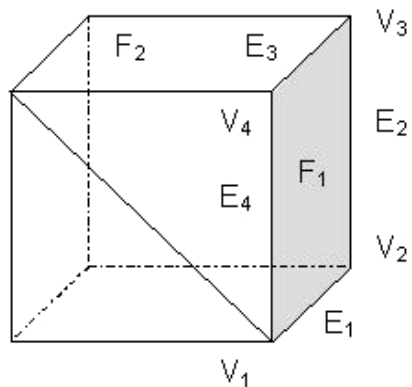
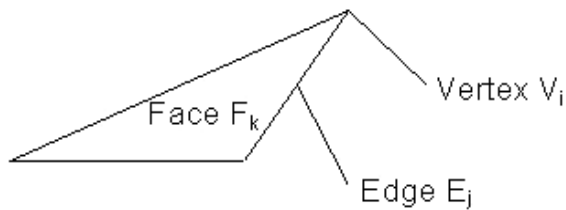
- Planar graph can be drawn on a plane with no intersection between edges
- Planar graph can be drawn (**embedded in**) on a sphere or any other genus 0 surface (without through holes) with no edge intersections
- Any graph can be embedded in some surface with **genus  $G$**  (number of through holes)
- **BRep can be considered a graph embedded in a surface**



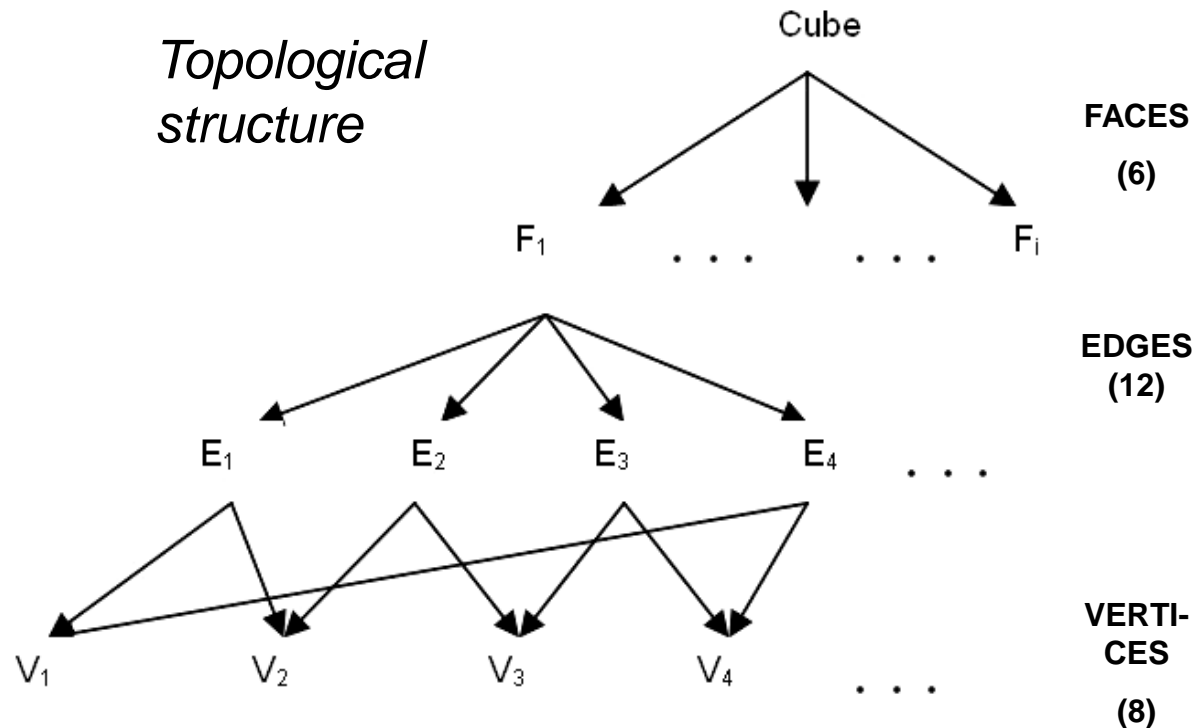


# BRep of a Cube

Example: A boundary representation for a cube



*Topological structure*





# Polyhedra and Euler's Formula

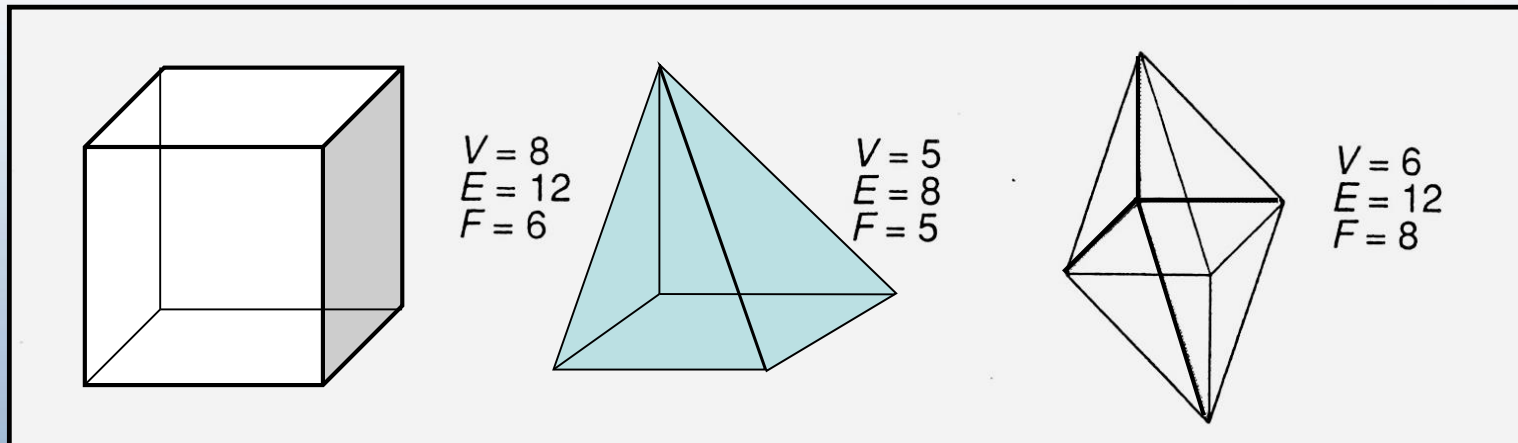
---

- A 3D polyhedron is a solid that is bounded by a set of polygons:
  - each edge connects two vertices and is shared by exactly two faces
  - at least three edges meet at each vertex
  - faces do not interpenetrate.



# Polyhedra and Euler's Formula

- A simple polyhedron can be deformed into a sphere (no through holes).



Some simple polyhedron with their V, E and F values.



- The BRep of simple polyhedron satisfies Euler's formula:

$$V - E + F = 2$$

where






V is the number of vertices

E is the number of edges

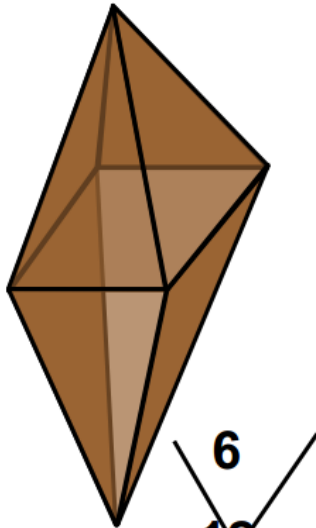
F is the number of faces.



# Polyhedra and Euler's Formula

Name	Image	Vertices $V$	Edges $E$	Faces $F$	Euler characteristic: $V - E + F$
Tetrahedron		4	6	4	2
Hexahedron or cube		8	12	6	2
Octahedron		6	12	8	2
Dodecahedron		20	30	12	2
Icosahedron		12	30	20	2

# Polyhedra and Euler's Formula



~~6  
12  
9~~

$V = 6$   
 $E = 12$   
 $F = 8$

$$V - E + F = 2$$

where

V is the number of vertices

E is the number of edges

F is the number of faces.

$$6 - 12 + 8 = 2$$

The internal face has to be removed as it violates the requirement that each edge is shared by exactly two faces.



# Generalized Euler's Formula

- The BRep of solids that have faces with holes satisfies the generalized Euler's formula:

$$\mathbf{V - E + F - H = 2 ( C - G )}$$

where

V is the number of vertices

E is the number of edges

F is the number of faces

H is the number of holes in the faces

C is the number of separate components (parts)

G is the genus – number of holes passing through the solid (for sphere  $G=0$ , torus  $G = 1$ )

If  $C > 1$ , G is the sum of the genera of all its components.



# Generalized Euler's Formula

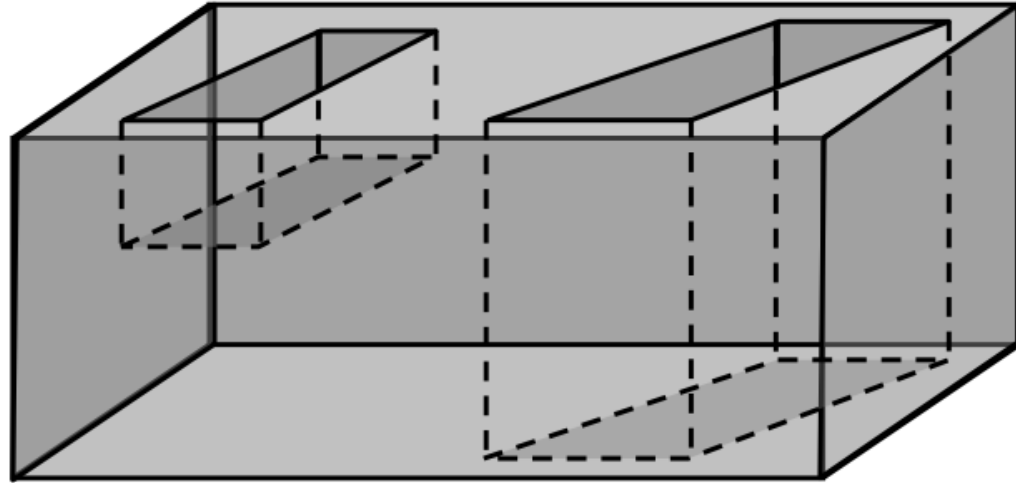
$$V = 24$$

$$E = 36$$

$$F = 15$$

$$H = 3$$

$$C = G = 1$$



$$V - E + F - H = 2(C - G)$$

V is the number of vertices

E is the number of edges

F is the number of faces

H is the number of holes in the faces

C is the number of separate components

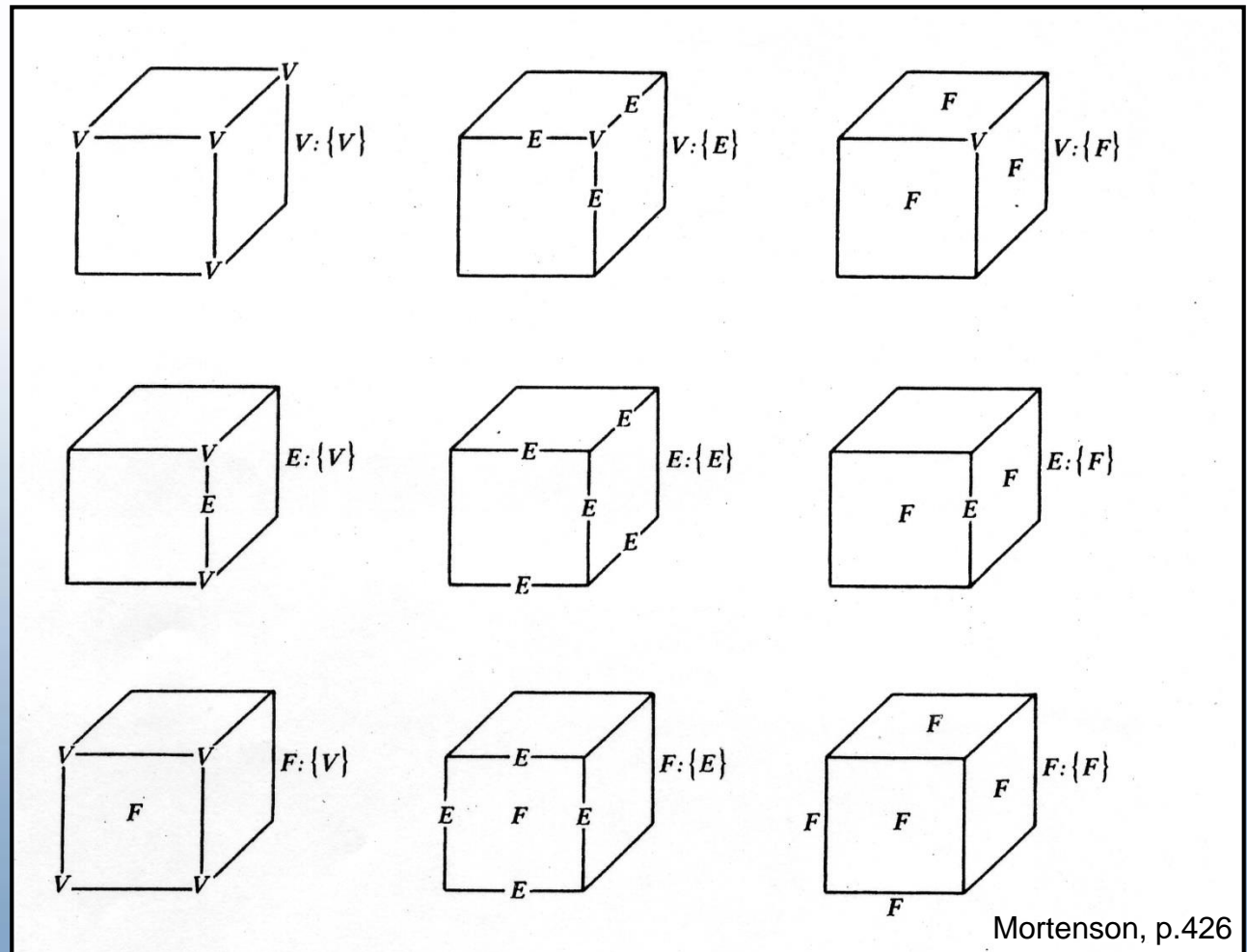
G is the genus





# Topological relationships in BRep

A polyhedron has **nine** classes of topological relationships between pairs of elements: vertices, edges, and faces.





Different applications need different adjacency information:

- $V: \{V\}$ ,  $E: \{V\}$ ,  $F: \{V\}$  in wireframe (vector) graphics to know how vertices are joined
- $V: \{F\}$  in set operations to know the ring of faces around the vertex
- $F: \{F\}$  adjacency among faces is needed in Euler operators.



# Winged-edge structure

---

This data structure represents the boundary of a manifold polyhedral object. The topological information is as follows:

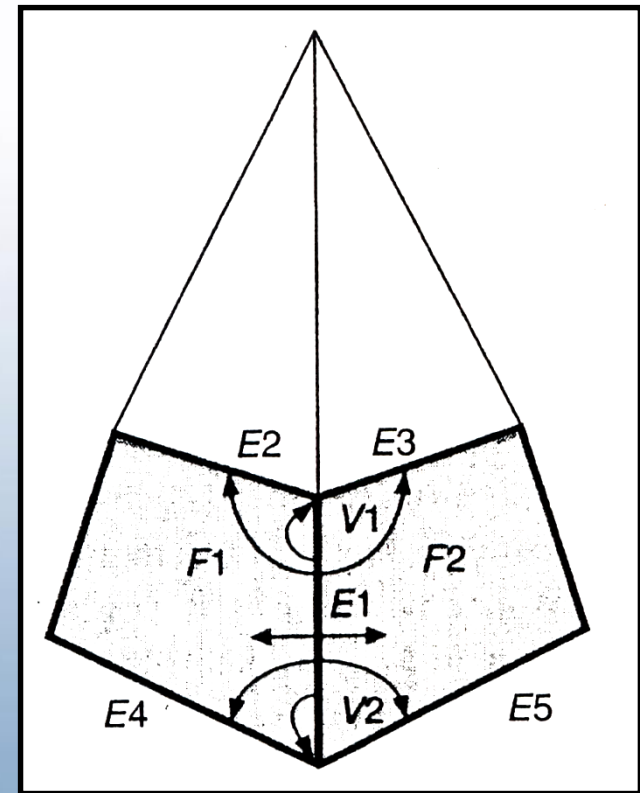
- Each face is bounded by a set of disjoint edge cycles. One cycle is the outside boundary of the face, the others bounding holes.
- Each vertex is adjacent to a circularly ordered set of edges, so the vertex table specifies one of these edges for each vertex.

# Winged-edge structure



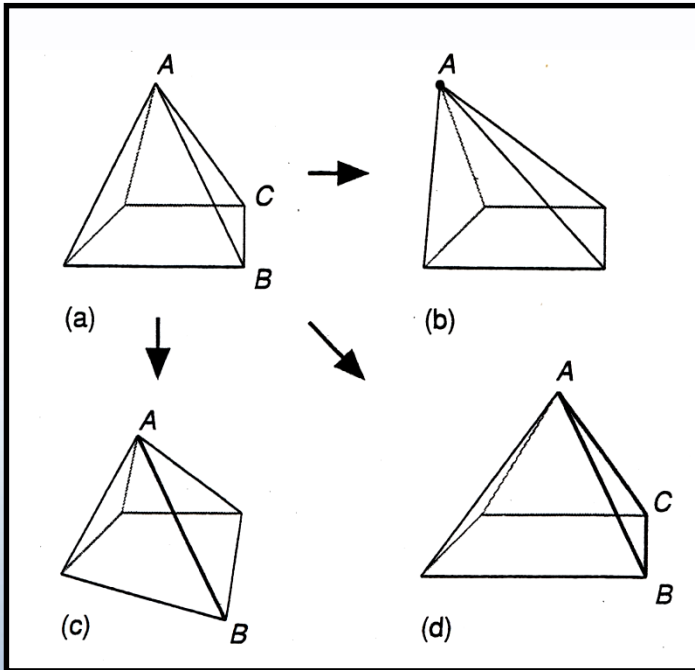
- For each edge the following information is given (see figure):
  - 1) Incident vertices ( $V1$ ,  $V2$ )
  - 2) Left and right adjacent face ( $F2$ ,  $F1$ )
  - 3) Two edges that share  $V1$  ( $E2$ ,  $E3$ )
  - 4) Two edges that share  $V2$  ( $E4$ ,  $E5$ )

This structure makes it possible to determine **in constant time** which vertices or faces are associated with an edge.

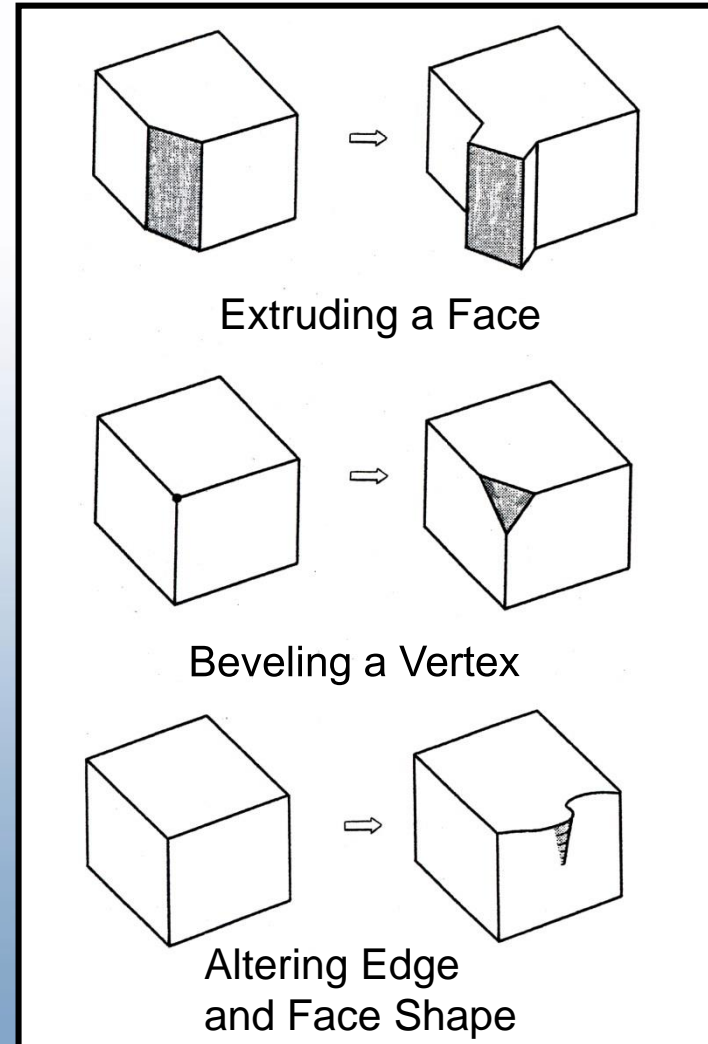




# Local modifications



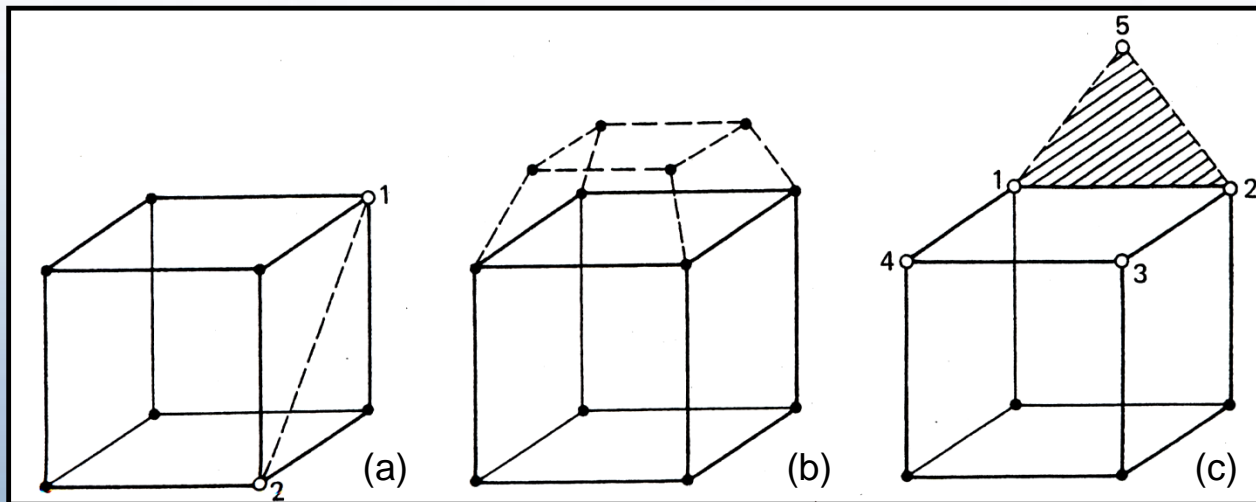
(a) An object on which tweaking operations are performed to move, (b) vertex A, (c) edge AB, (d) face ABC





# Euler operators

Euler operators transform the objects satisfying Euler's formula by adding and removing vertices, edges and faces.



(a, b) a cubical polyhedron is correctly modified;  
(c) the result is not a polyhedron because edges (1,5) and (2,5) are not shared by two faces each.



## Euler operators

A linear combination of five primitive operators (with their inverses) can represent all objects satisfying Euler formula:

1. Make (kill) an edge and a vertex (***mev / kev***)
2. Make a face and an edge (***mfe / kfe***)
3. Make a body, a face, and a vertex (***mbfv / kbfv***);
4. Make a cavity, or passage, and a body (***mrbb / krb***);
5. Make an edge and kill a hole (***me-kh***).



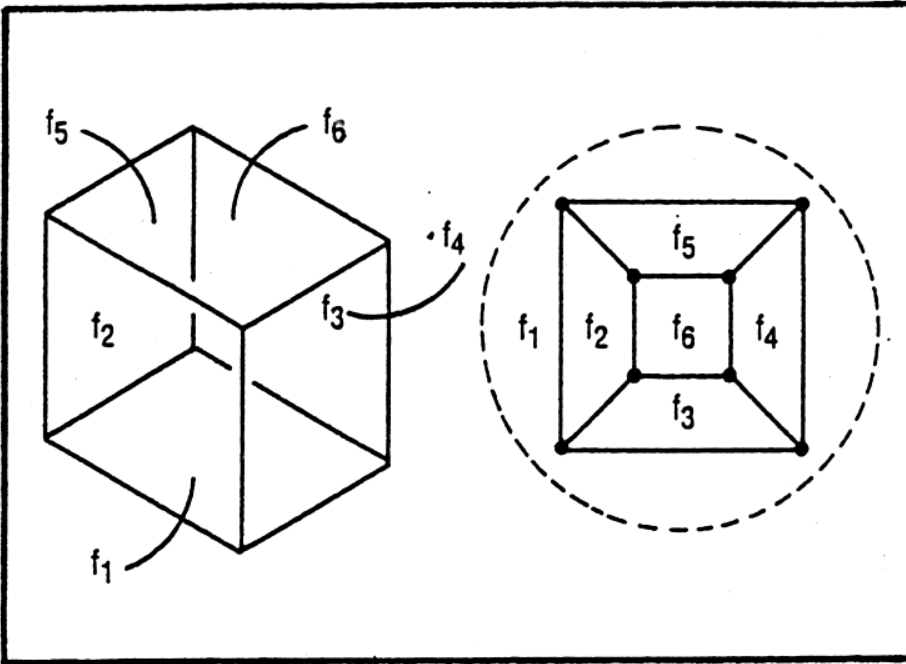
## Advantages of Euler operators:

- Ensured topological validity of the resulting solids
- Intermediate language isolating high-level operations from the underlying data structures





# Euler operators in the GWB system



## Euler operators

OPERATOR	EXPLANATION
$mvsf(f, v)$	MAKE VERTEX, SOLID, FACE
$kvsf()$	KILL VERTEX, SOLID, FACE
$mev(v_1, v_2, e)$	MAKE EDGE, VERTEX
$kev(e, v)$	KILL EDGE, VERTEX
$mef(v_1, v_2, f_1, f_2, e)$	MAKE EDGE, FACE
$kef(e)$	KILL EDGE, FACE
$kemr(e)$	KILL EDGE, MAKE RING
$mekr(v_1, v_2, e)$	MAKE EDGE, KILL RING
$kfmrh(f_1, f_2)$	KILL FACE, MAKE RING, HOLE
$mfkrh(f_1, f_2)$	MAKE FACE, KILL RING, HOLE
$semv(e_1, v, e_2)$	SPLIT EDGE, MAKE VERTEX
$jekv(e_1, e_2)$	JOIN EDGES, KILL VERTEX

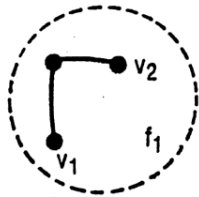
Cube topology: a plane model



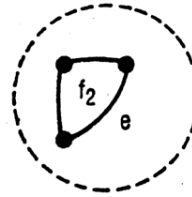
mvsf  
 ← kvsf  
 (a)



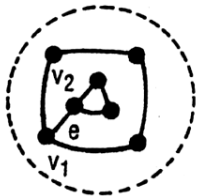
mev  
 ← kev  
 (b)



mef  
 ← kef  
 (c)



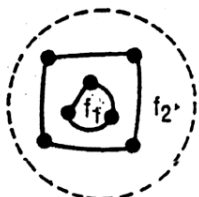
mef  
 ← kef  
 (d)



kemr  
 ← mekr  
 (e)



kemr  
 ← mekr  
 (f)



kfmrh  
 ← mfkrrh  
 (g)



kfmrh  
 ← mfkrrh  
 (h)



semv  
 ← jekv  
 (i)



semv  
 ← jekv  
 (j)

# Euler operators in the GWB system



# Contents

---

- Boundary Representation BRep
- Euler's formula
- Winged-edge structure
- Euler operators
- **CSG tree**
- Point-membership classification algorithm
- References



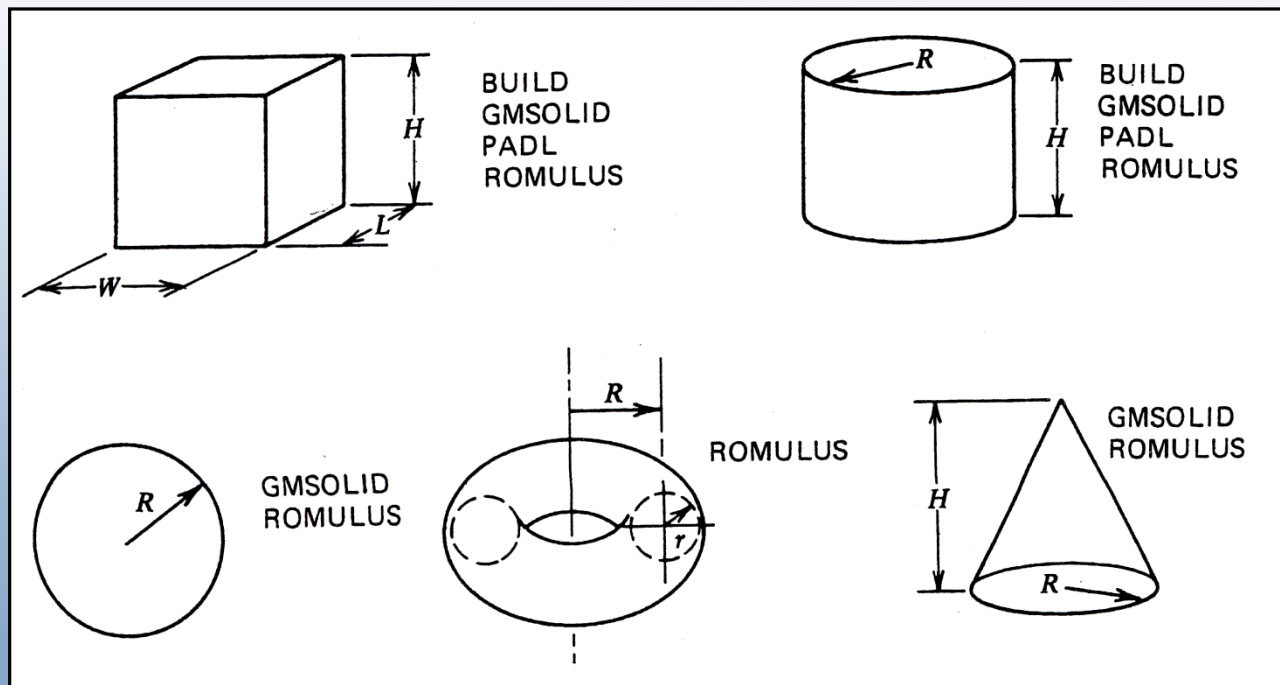
# Constructive Solid Geometry

## CSG: primitives

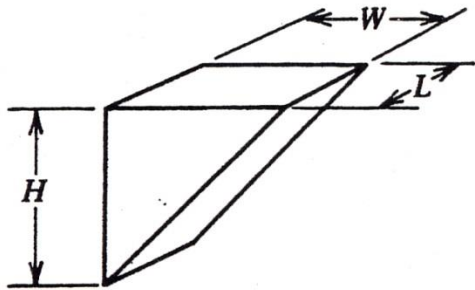
In CSG, simple primitives are combined by means of regularized set operations and rigid motions.

Standard CSG primitives:

**block, cylinder, sphere, cone, and torus**

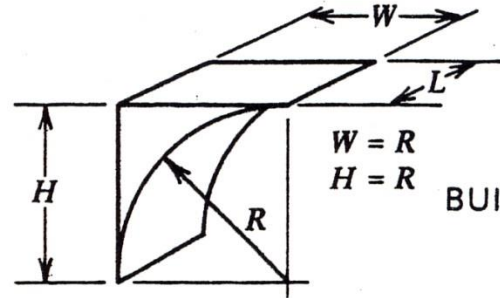


# CSG representation: primitives



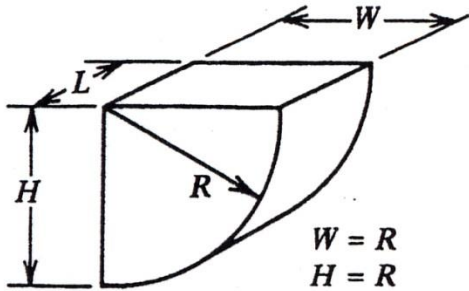
BUILD

(c)

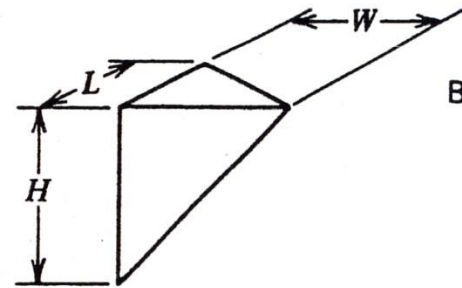


BUILD

(d)



BUILD

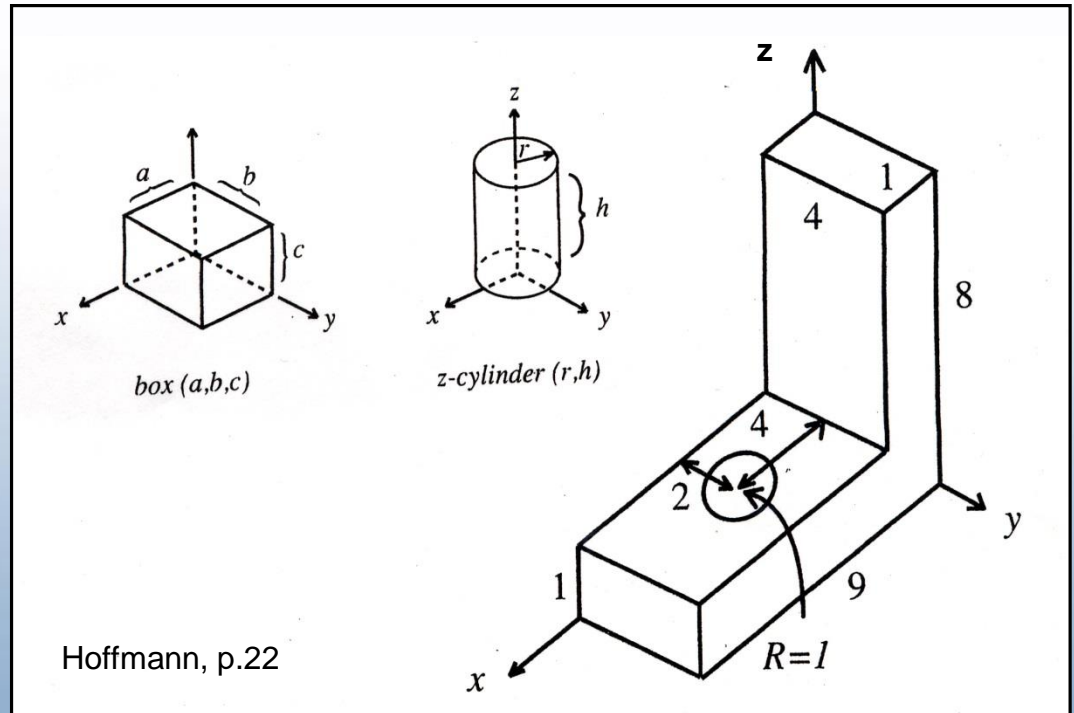


BUILD

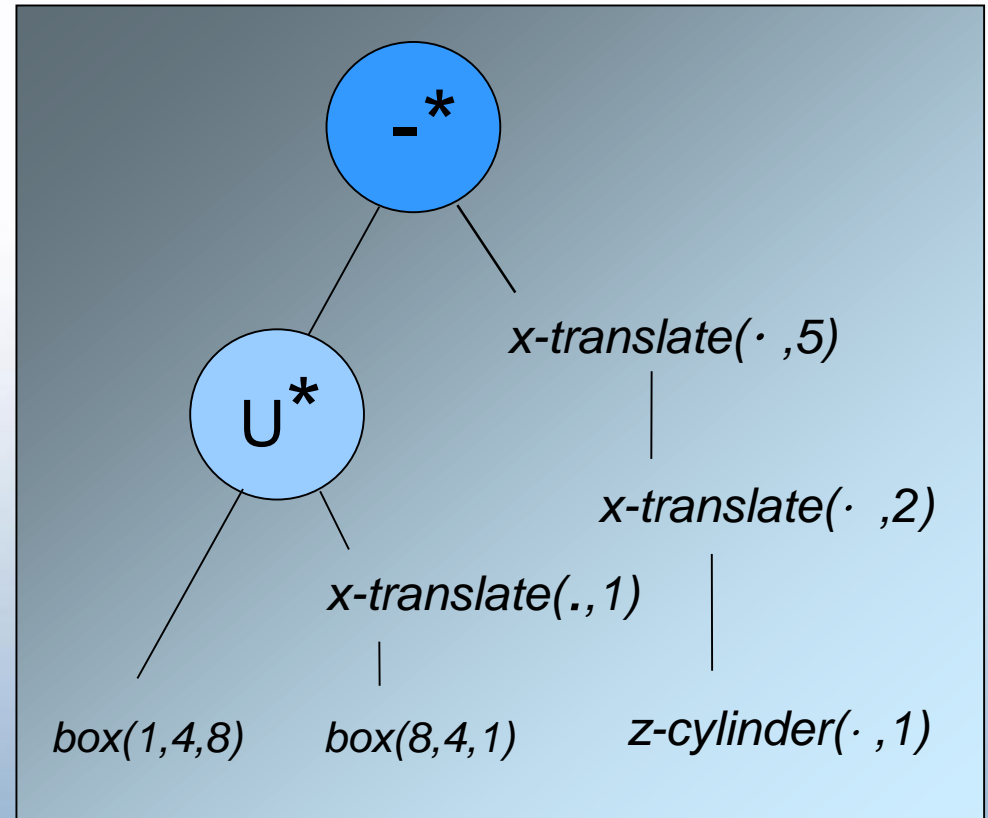
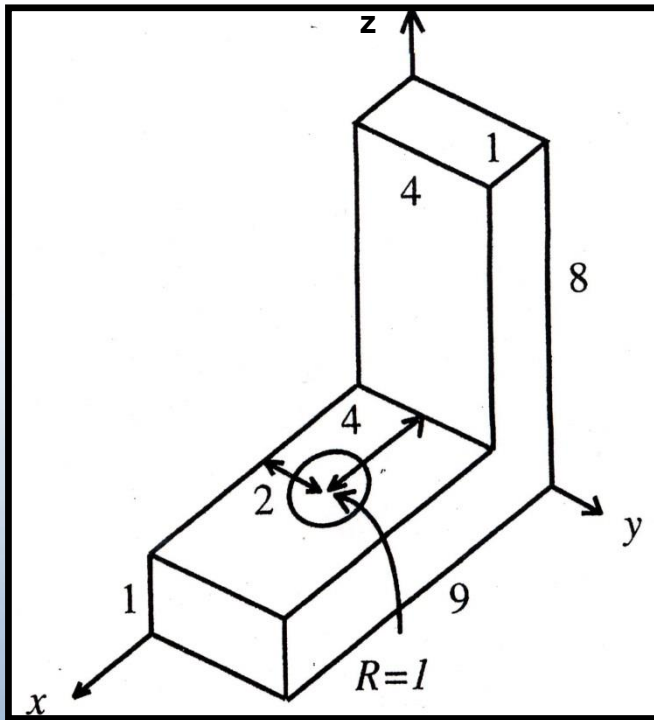


# CSG Tree

- An object is represented as a **binary tree** with operations at the internal nodes and primitives at the leaves
- Nodes: regularized set operations or rigid motions (translation, rotation)



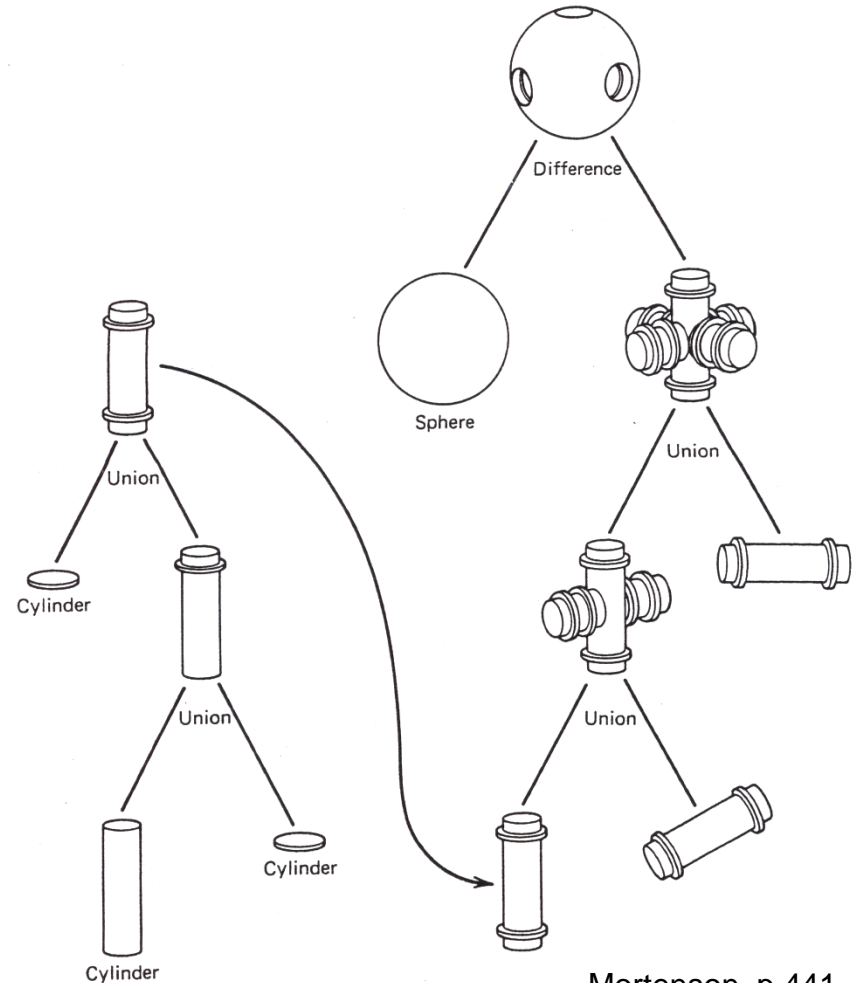
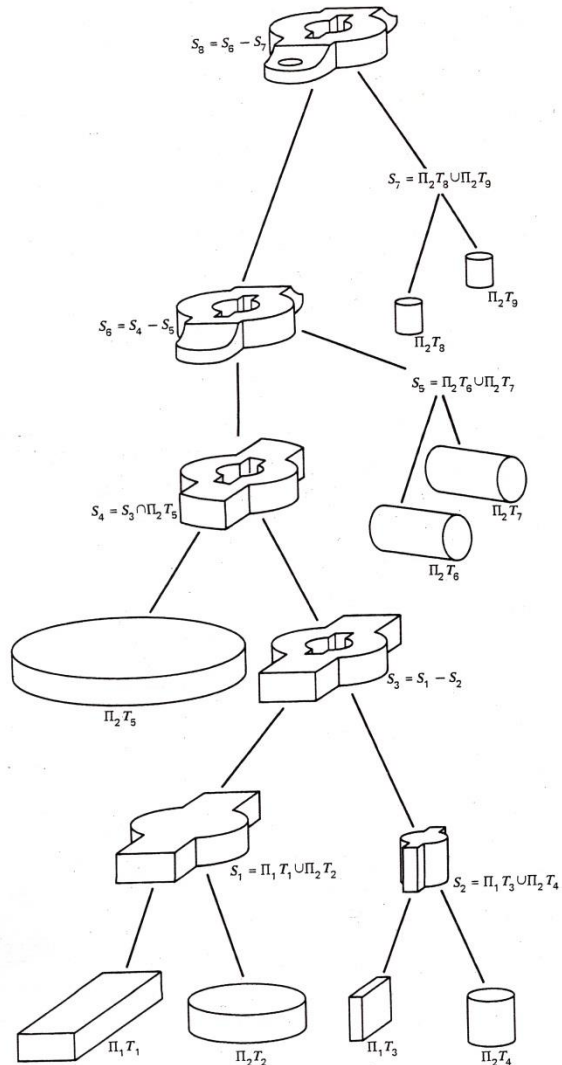
# CSG Tree



$(\text{block}(1,4,8) \cup \text{x-translate}(\text{block}(8,4,1),1) - \text{x-translate}(\text{y-translate}(\text{z-cylinder}(1,1),2),5))$



# CSG Tree: Examples







# Point membership classification

---

A solid  $S$  and a point  $X$  are given.

Query (binary predicate):

**Is  $X$  inside or outside of  $S$  ?**

Point membership classification (PMC) function

$$M[X,S] = (XinS, XoutS)$$

Query (ternary predicate):

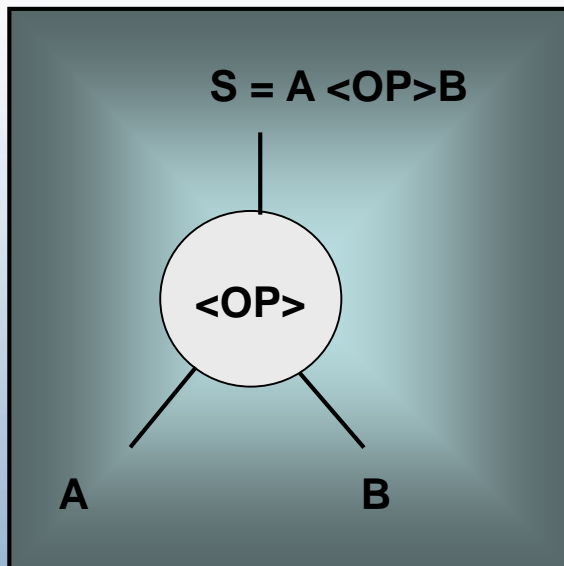
**Is  $X$  inside, outside or on the boundary of  $S$  ?**

Point membership classification (PMC) function

$$M[X,S] = (XinS, XonS, XoutS)$$



## Divide-and-conquer paradigm



$F(S) \leftarrow$  IF (S is a primitive)  
THEN prim - f(S) ELSE combine  
( f(A), f(B), <OP> )

The divide and conquer paradigm.  
<OP> is a regularized operator  
(  $\cup^*$  ,  $\cap^*$  , or  $-^*$  )



```
M [ X, S ] ← IF (S is a primitive)  
  THEN prim-M (X,S)  
  ELSE combine ( M [X, left-subtree(S),  
    M [ X, right-subtree (S)], root <S>))
```

- *prim-M* is a primitive classification procedure and must produce “in”, “on” or “out” answers for a given point
- *combine* applies set operations (three-valued logic!) to its arguments
- **postorder tree traversal**



# Recursive PMC algorithm structure

---

1. Point coordinates  $(x,y,z)$  are sent to the root of the CSG tree.
2. Downward propagation
  - Coordinates are propagated into the tree down to the leaves, possibly altered at the motion nodes.
  - At each leaf, the final point coordinates describe the same point, but in the local coordinate frame of the primitive solid.



## Downward propagation cases:

- 1) If  $(x,y,z)$  arrives at a set-theoretic operation node, it is passed unchanged to the two subtrees
- 2) If  $(x,y,z)$  arrives at a motion node, the inverse transformation is applied to  $(x,y,z)$ , resulting in new local coordinates  $(x',y',z')$ , which are sent to the two subtrees
- 3) If  $(x,y,z)$  arrives at a leaf, the point is classified against the primitive, and the classification is returned to the parent of the leaf

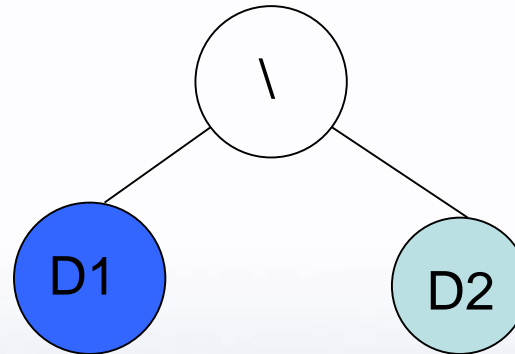
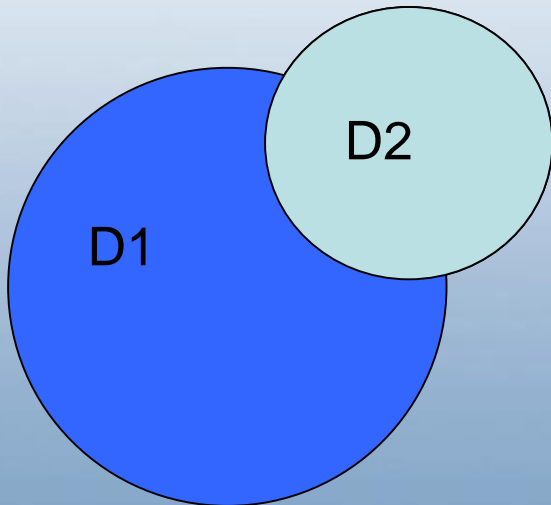
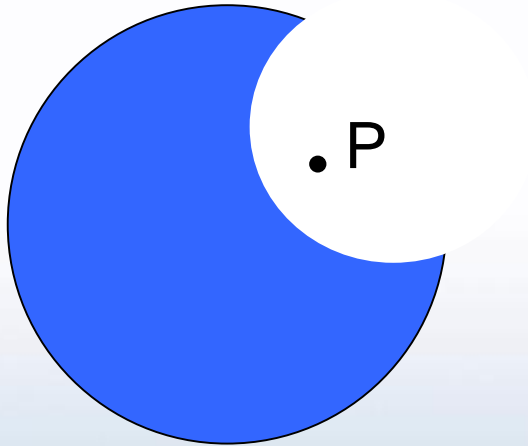


## 3. Upward propagation

- Classifications from the subtrees are combined in the set-theoretic operation nodes
- No work is done at motion nodes.



# PMC Example



$M [ X, S ] \leftarrow$  IF (S is a primitive)  
THEN prim-M (X,S)  
ELSE combine ( M [X, left-subtree(S),  
M [ X, right-subtree (S)], root <S>))

Postorder tree traversal:

D1-D2-\

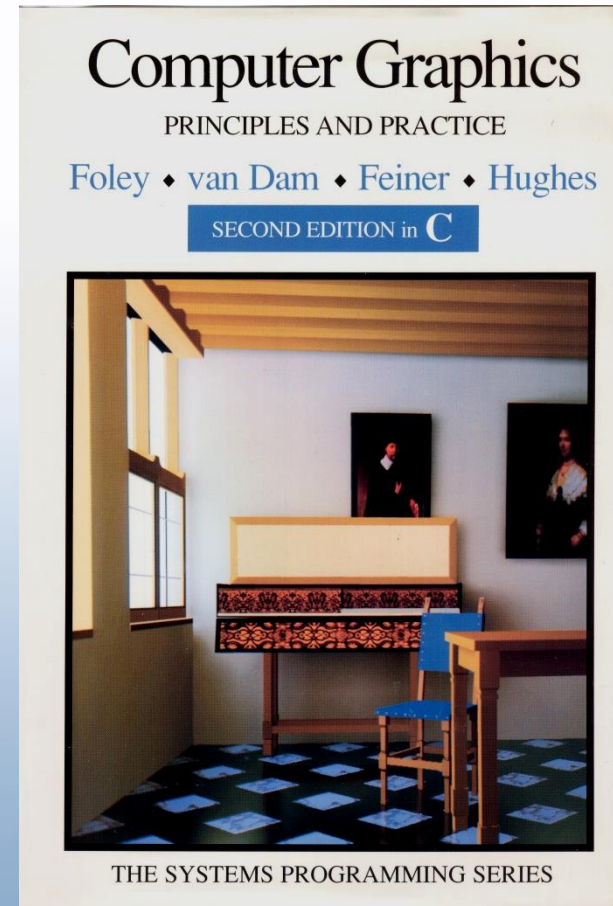
PMC steps:

- 1)
- 2)
- 3)



# References

- James D. Foley, Andries van Dam, Steven K. Feiner, John F. Hughes, Computer Graphics: Principles and Practice (2nd Edition in C ), Addison-Wesley, Reading, MA, 1997.







- Michael E. Mortenson, Geometric Modeling, John Wiley and Sons, 1985. Second edition, 1997. Third edition, 2006.
- Christoph M. Hoffmann Geometric and Solid Modeling. An Introduction, Morgan Kaufmann Publishers, 1989, 338 p.

